#### Minimization of FSA

Data Structures and Algorithms for Comp (ISCL-BA-07) nal Linguistics III

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### Finding equivalent states



Their right languages are the same

## Minimization by partitioning



- Accepting & non-accept partition
- If any two nodes go to different sets for any of the symbols split
- $Q_1 = \{0,3\}, Q_2 = \{1\}, Q_4 = \{2\}, Q_2 = \{4,5\}$
- Stop when we cannot split any of the sets, merge the indistinguishable states

## Minimization by partitioning

## Create a state-by-state table, mark distinguishers pairs: $(q_1,q_2)$ such that $(\Delta(q_1,x),\Delta(q_2,x))$ is a distinguishable pair for any $x\in \Sigma$



#### Minimization by partitioning



## Minimization by partitioning



Create a state-by-state table, mark distinguishable pairs:  $(q_1,q_2)$  such that  $(\Delta(q_1,x),\Delta(q_2,x))$  is a distinguishable pair for any  $x\in \Sigma$ 



#### DFA minimization

# \* For any regular language, there is a unique minimal DFA

- By finding the minimal DFA, we can also prove equivalence (or not) of different FSA and the languages they recognize
- In general the idea is:
- \* us goverfall the 1000 ts:

   Thereo area pure numericable states (casy)

   Merge equivalent states

  \* There are two well-known algorithms for minimization:

   Hoperoff's algorithm: find and elliminate equivalent states by partitioning the set of states

   Benzowowski's algorithm: 'double reversal'

## Finding equivalent states



### Minimization by partitioning



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## Minimization by partitioning



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Minimization by partitioning

Create a state-by-state table, mark distinguishable pairs: (q<sub>1</sub>, q<sub>2</sub>) such that (Δ(q<sub>1</sub>, x), Δ(q<sub>2</sub>, x)) is a distinguishable pair for any x ∈ Σ



#### Minimization by partitioning

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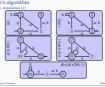


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- The algorithm can be cell to visit carefully

Brzozowski's algorithm



An exercise



Minimization algorithms

- There are many versions of the 'partitioning' algorithm. General idea is to form equivalence classes based on right-language of each state.

  Partitioning algorithm has O(n log n) complexity
- 'Double reversal' algorithm has exponential worst-time complexity
- Double reversal algorithm can also be used with NFAs (resulting in the minimal equivalent DFA NFA minimization is intractable)
- ner, different alg
- different input
- Reading suggesti ion: Hopcroft and Ullman (1979, Ch. 2&3), Jurafsky and Martin (2009, Ch. 2)
- FST
  - FSA and regular language

Acknowledgments, credits, references

- Hopcroft, John E. and Jeffrey D. Ullman (1979). Introduction to Automata Theory, Languages, and Computation. Addison-Wesley Series in Computer Science and
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